

LECTURE - XVIII
CONCURRENT PROGRAMMING

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Roadmap

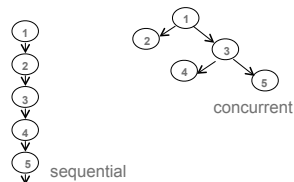
- Concurrent Programming
 - Shared Memory vs Message Passing
 - Divide and Conquer
 - Threads vs Processes
 - POSIX Threads



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Concurrent Programming

- So far, we have focused on **sequential programming**: all computational tasks are executed in sequence, one after the other.
- Next three lectures, we will focus on **concurrent programming**: multiple computational tasks are executed simultaneously, at the same time.



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Concurrent Programming

- Implementation of concurrent tasks:
 - as separate programs
 - as a set of processes or threads created by a single program
- Execution of concurrent tasks:
 - on a single processor
 - **Multithreaded programming**
 - on several processors in close proximity
 - **Parallel computing**
 - on several processors distributed across a network
 - **Distributed computing**

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Communication Between Tasks

Interaction or communication between concurrent tasks can be done via:

- **Shared memory:**
 - all tasks have access to the same physical memory
 - they can communicate by altering the contents of shared memory
- **Message passing:**
 - no common/shared physical memory
 - tasks communicate by exchanging messages

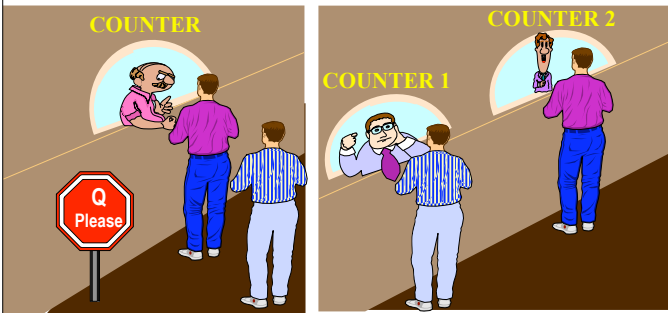
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Motivation

- Increase the performance by running more than one task at a time.
 - divide the program into n smaller pieces, and run it n times faster using n processors
- To cope with independent physical devices.
 - do not wait for a blocked device, perform other operations at the background

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Serial vs Parallel



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Divide and Compute

$$x1 + x2 + x3 + x4 + x5 + x6 + x7 + x8$$

How many operations with sequential programming?
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Step 1: $x1 + x2$

Step 2: $x1 + x2 + x3$

Step 3: $x1 + x2 + x3 + x4$

Step 4: $x1 + x2 + x3 + x4 + x5$

Step 5: $x1 + x2 + x3 + x4 + x5 + x6$

Step 6: $x1 + x2 + x3 + x4 + x5 + x6 + x7$

Step 7: $x1 + x2 + x3 + x4 + x5 + x6 + x7 + x8$

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Divide and Compute

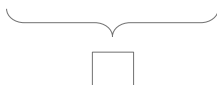
$$x1 + x2 + x3 + x4 + x5 + x6 + x7 + x8$$



Step 1: parallelism = 4



Step 2: parallelism = 2



Step 3: parallelism = 1

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Gain from parallelism

In theory:

- dividing a program into n smaller parts and running on n processors results in n time speedup

In practice:

- This is not true, due to
 - Communication costs
 - Dependencies between different program parts
 - Eg. the addition example can run only in $\log(n)$ time not $1/n$

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Prevent Blocking

- Do not wait for a blocked device, perform other operations at the background
 - During I/O perform computation
 - During continuous visualization, handle key strokes and I/O
 - Eg. video games
 - While listening to network, perform other operations
 - Listening to multiple sockets at the same time
 - Concurrent I/O, concurrent transfers
 - Eg. Web browsers

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Threads vs Processes

Process Spawning:

Process creation involves the following four main actions:

- setting up the process control block,
- allocation of an address space and
- loading the program into the allocated address space and
- passing on the process control block to the scheduler

Thread Spawning:

- Threads are created *within and belonging to* processes
- All the threads created within one process share the resources of the process including the address space
- Scheduling is performed on a per-thread basis.
- The thread model is a *finer grain scheduling model* than the process model
- Threads have a similar *lifecycle* as the processes and will be managed mainly in the same way as processes are

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Threads vs Processes

- Heavyweight Process = Process
- Lightweight Process = Thread

Advantages (Thread vs. Process):

- Much quicker to create a thread than a process
- Much quicker to switch between threads than to switch between processes
- Threads share data easily

Disadvantages (Thread vs. Process):

- Processes are more flexible
 - They don't have to run on the same processor
- No security between threads: One thread can stomp on another thread's data
- For threads which are supported by user thread package instead of the kernel:
 - If one thread blocks, all threads in task block.

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Synchronization

- Mechanism that allows the programmer to control the relative order in which operations occur in different threads or processes.

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Synchronization - Threads

Int sum = 0;

Thread 1:

```
int t;
lock(sum);
sum = sum + x;
t = sum;
....
unlock(sum);
```

Thread 2:

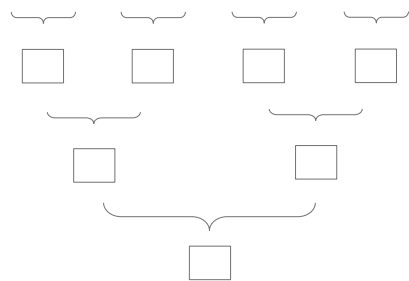
```
int t;
lock(sum);
sum = sum + y;
t = sum;
...
unlock(sum);
```

Use of **semaphores** for thread synchronization!

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Synchronization - Processes

$x1 + x2 + x3 + x4 + x5 + x6 + x7 + x8$



Step 1: parallelism = 4

Step 2: parallelism = 2

Step 3: parallelism = 1

Wait for a message from other processes before continuing processing!

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On a single processor machine

- You can have multiple threads
- You can also have multiple processes and have the effect of concurrency
 - timesharing

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Thread Creation

• pthread_create

```
// creates a new thread executing start_routine
int pthread_create(pthread_t *thread,
                  const pthread_attr_t *attr,
                  void *(*start_routine)(void*), void *arg);
```

• pthread_join

```
// suspends execution of the calling thread until the target
// thread terminates
int pthread_join(pthread_t thread, void **value_ptr);
```

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Mutual Exclusion

- **pthread_mutex_lock**
// blocks until mutex is available, and then locks it
int pthread_mutex_lock(pthread_mutex_t *mutex);

pthread_mutex_unlock
// unlocks the mutex
int pthread_mutex_unlock(pthread_mutex_t *mutex);

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Thread Example

```
int main()
{
    pthread_t thread1, thread2; /* thread variables */

    pthread_create (&thread1, NULL, (void *) &print_message_function,
                  (void*)"hello ");
    pthread_create (&thread2, NULL, (void *) &print_message_function,
                  (void*)"world!\n");

    pthread_join(thread1, NULL);
    pthread_join(thread2, NULL);

    exit(0);
}
```

Why use pthread_join?

To force main block to wait for both threads to terminate, before it exits.
If main block exits, both threads exit, even if the threads have not finished their work.

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Thread Example (cont.)

```
void print_message_function ( void *ptr )
{
    char *cp = (char*)ptr;

    for (i=0;i<NUM;i++){
        printf("%s ", cp);
        fflush(stdout);
    }

    pthread_exit(0); /* exit */
}
```

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Interthread Cooperation

```
void* print_count ( void *ptr );
void* increment_count ( void *ptr );

int NUM=5;
int counter =0;

int main()
{
    pthread_t thread1, thread2;

    pthread_create (&thread1, NULL, increment_count, NULL);
    pthread_create (&thread2, NULL, print_count, NULL);

    pthread_join(thread1, NULL);
    pthread_join(thread2, NULL);

    exit(0);
}
```

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Interthread Cooperation (cont.)

```
void* print_count ( void *ptr )
{
    int i;
    for (i=0;i<NUM;i++){
        printf("counter = %d\n", counter);
        sleep(1);
    }
    pthread_exit(0);
}

void* increment_count ( void *ptr )
{
    int i;
    for (i=0;i<NUM;i++){
        counter++;
        sleep(1);
    }
    pthread_exit(0);
}
```

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2-Thread Word Counter

```
int total_words;

main(int ac, char *av[])
{
    pthread_t t1, t2; /* two threads */
    void *count_words(void *);

    if ( ac != 3 ){
        printf("usage: %s file1 file2\n", av[0]);
        exit(1);
    }
    total_words = 0;

    pthread_create(&t1, NULL, count_words, (void *) av[1]);
    pthread_create(&t2, NULL, count_words, (void *) av[2]);
    pthread_join(t1, NULL);
    pthread_join(t2, NULL);
    printf("%5d: total words\n", total_words);
}
```

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2-Thread Word Counter (cont.)

```
void *count_words(void *f)
{
    char *filename = (char *) f;
    FILE *fp;
    int c, prevc = '\0';

    if ( (fp = fopen(filename, "r")) != NULL ){
        while( ( c = getc(fp)) != EOF ){
            if ( !isalnum(c) && !isalnum(prevc) )
                total_words++;
            prevc = c;
        }
        fclose(fp);
    } else
        perror(filename);
    return NULL;
}
```

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2-Thread Word Counter, Mutex

```
pthread_mutex_t counter_lock = PTHREAD_MUTEX_INITIALIZER;

int total_words ; /* the counter and its lock */

main(int ac, char *av[])
{
    pthread_t t1, t2; /* two threads */
    void *count_words(void *);

    if ( ac != 3 ){
        printf("usage: %s file1 file2\n", av[0]);
        exit(1);
    }
    total_words = 0;

    pthread_create(&t1, NULL, count_words, (void *) av[1]);
    pthread_create(&t2, NULL, count_words, (void *) av[2]);
    pthread_join(t1, NULL);
    pthread_join(t2, NULL);
    printf("%5d: total words\n", total_words);
}
```

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2-Thread Word Counter, Mutex (cont.)

```
void *count_words(void *f)
{
    char *filename = (char *) f;
    FILE *fp;
    int c, prevc = '\0';

    if ( (fp = fopen(filename, "r")) != NULL ){
        while( ( c = getc(fp)) != EOF ){
            if ( !isalnum(c) && !isalnum(prevc) ){
                pthread_mutex_lock(&counter_lock);
                total_words++;
                pthread_mutex_unlock(&counter_lock);
            }
            prevc = c;
        }
        fclose(fp);
    } else
        perror(filename);
    return NULL;
}
```

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2-Thread Word Counter, Arg Pass

```
main(int ac, char *av[])
{
    pthread_t t1, t2; /* two threads */
    struct arg_set args1, args2; /* two argsets */
    void *count_words(void *);

    if ( ac != 3 ){
        printf("usage: %s file1 file2\n", av[0]);
        exit(1);
    }
    args1.fname = av[1];
    args1.count = 0;
    pthread_create(&t1, NULL, count_words, (void *) &args1);

    args2.fname = av[2];
    args2.count = 0;
    pthread_create(&t2, NULL, count_words, (void *) &args2);

    pthread_join(t1, NULL);
    pthread_join(t2, NULL);
    printf("%5d: %s\n", args1.count, av[1]);
    printf("%5d: %s\n", args2.count, av[2]);
    printf("%5d: total words\n", args1.count+args2.count);
}
```

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2-Thread Word Counter, Arg Pass (cont.)

```
struct arg_set { /* two values in one arg */
    char *fname; /* file to examine */
    int count; /* number of words */
};

void *count_words(void *a)
{
    struct arg_set *args = a; /* cast arg back to correct type */
    FILE *fp;
    int c, prevc = '\0';

    if ( (fp = fopen(args->fname, "r")) != NULL ){
        while( ( c = getc(fp)) != EOF ){
            if ( !isalnum(c) && !isalnum(prevc) )
                args->count++;
            prevc = c;
        }
        fclose(fp);
    } else
        perror(args->fname);
    return NULL;
}
```

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