

CSC 4103 - Operating Systems  
Fall 2009

LECTURE - XX  
FILE SYSTEMS

Tevfik Koşar

Louisiana State University  
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## File Systems

- Provides organized and efficient access to data on secondary storage:
  1. Organizing data into files and directories and supporting primitives to manipulate them (create, delete, read, write etc)
  2. Improve I/O efficiency between disk and memory (perform I/O in units of blocks rather than bytes)
  3. Ensure confidentiality and integrity of data
- Contains file structure via a File Control Block (FCB)
  - Ownership, permissions, location..

## A Typical File Control Block

file permissions
file dates (create, access, write)
file owner, group, ACL
file size
file data blocks or pointers to file data blocks

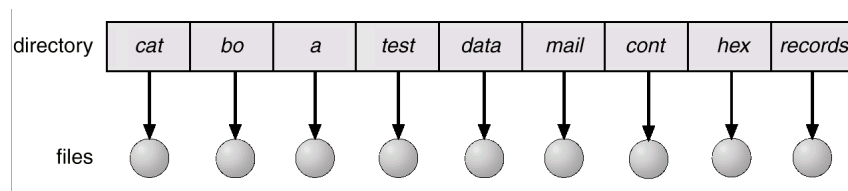
## Directories

- Directories are special files that keep track of other files
  - ✓ the collection of files is systematically organized
  - ✓ first, disks are split into partitions that create logical volumes (can be thought of as “virtual disks”)
  - ✓ second, each partition contains information about the files within
  - ✓ this information is kept in entries in a **device directory** (or volume table of contents)
  - ✓ the directory is a symbol table that translates file names into their entries in the directory
    - it has a logical structure
    - it has an implementation structure (linked list, table, etc.)

## Directories

### ➤ Single-level directory structure

- ✓ simplest form of logical organization: one global or **root** directory containing all the files
- ✓ problems
  - global namespace: unpractical in multiuser systems
  - no systematic organization, no groups or logical categories of files that belong together



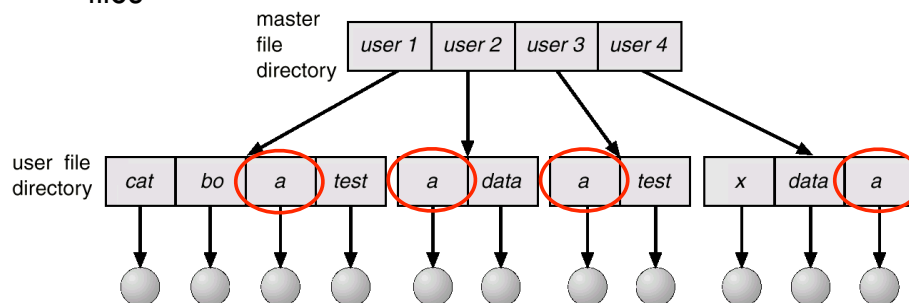
Single-level directory

Silberschatz, A., Galvin, P. B. and Gagne, G. (2003)  
Operating Systems Concepts with Java (6th Edition).

## Directories

### ➤ Two-level directory structure

- ✓ in multiuser systems, the next step is to give each user their own private directory
- ✓ avoids filename confusion
- ✓ however, still no grouping: not satisfactory for users with many files

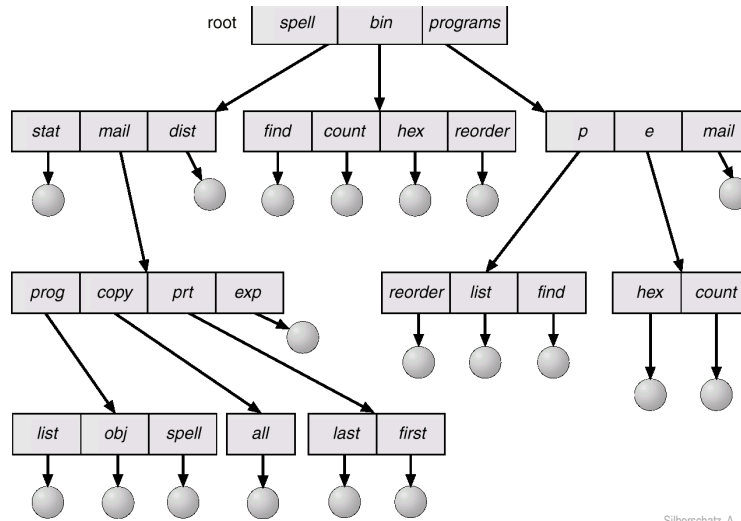


Two-level directory

Silberschatz, A., Galvin, P. B. and Gagne, G. (2003)  
Operating Systems Concepts with Java (6th Edition).

## Directories

### ➤ Tree-structured directory structure



Tree-structured directory

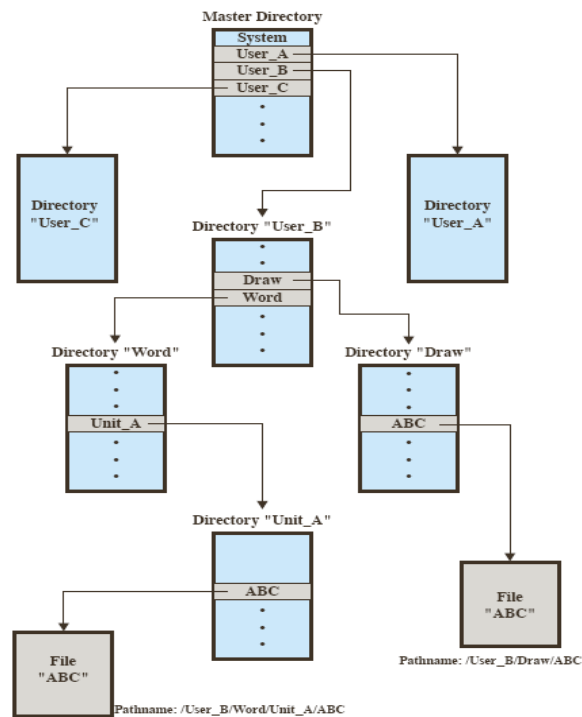
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## Directories

### ➤ Tree-structured directory structure

- ✓ natural extension of the two-level scheme
- ✓ provides a general hierarchy, in which files can be grouped in natural ways
- ✓ good match with human cognitive organization: tendency to categorize objects in embedded sets and subsets
- ✓ navigation through the tree relies on **pathnames**
  - absolute pathnames start from the root, example: /jsmith/academic/teaching/cs446/assignment4/grades
  - relative pathnames start at from a current **working directory**, example: assignment4/grades
  - the current and parent directory are referred to as . and ..

# Directory Implementation



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# Directory Implementation

- **Linear list** of file names with pointer to the data blocks.
  - simple to program
  - time-consuming to execute
- **Hash Table** - linear list with hash data structure.
  - decreases directory search time
  - **collisions** - situations where two file names hash to the same location
  - fixed size

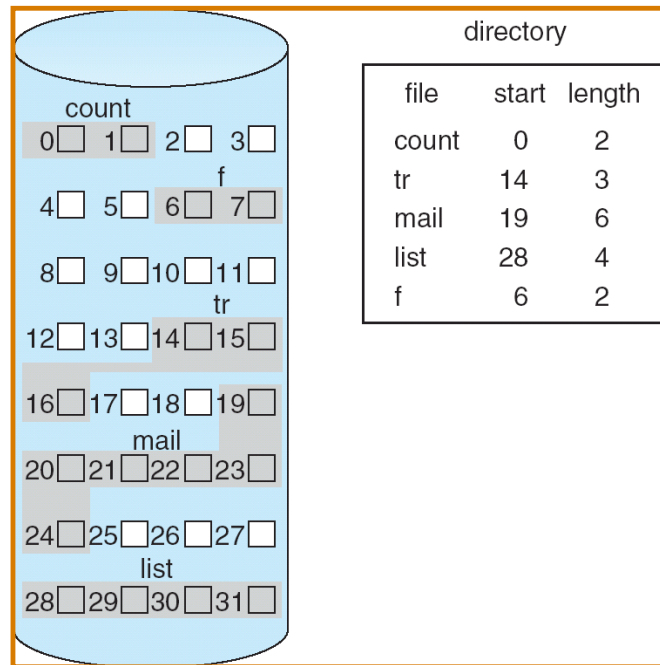
## Allocation Methods

- An allocation method refers to how disk blocks are allocated for files:
- **Contiguous allocation**
- **Linked allocation**
- **Indexed allocation**

## Contiguous Allocation

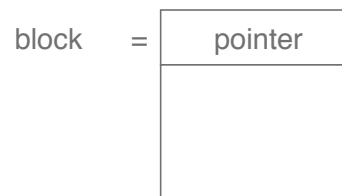
- Each file occupies a set of contiguous blocks on the disk
- + Simple - only starting location (block #) and length (number of blocks) are required
- - Wasteful of space (dynamic storage-allocation problem - fragmentation)
- - Files cannot grow

## Contiguous Allocation of Disk Space



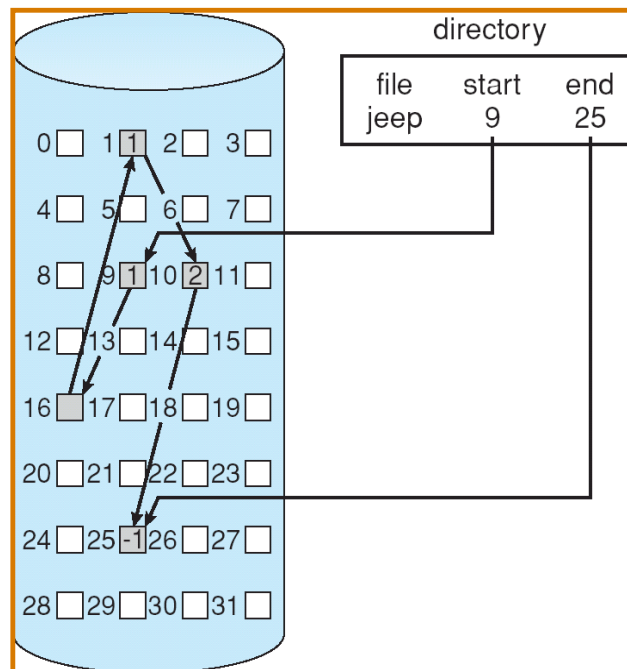
## Linked Allocation

- Each file is a linked list of disk blocks: blocks may be scattered anywhere on the disk.

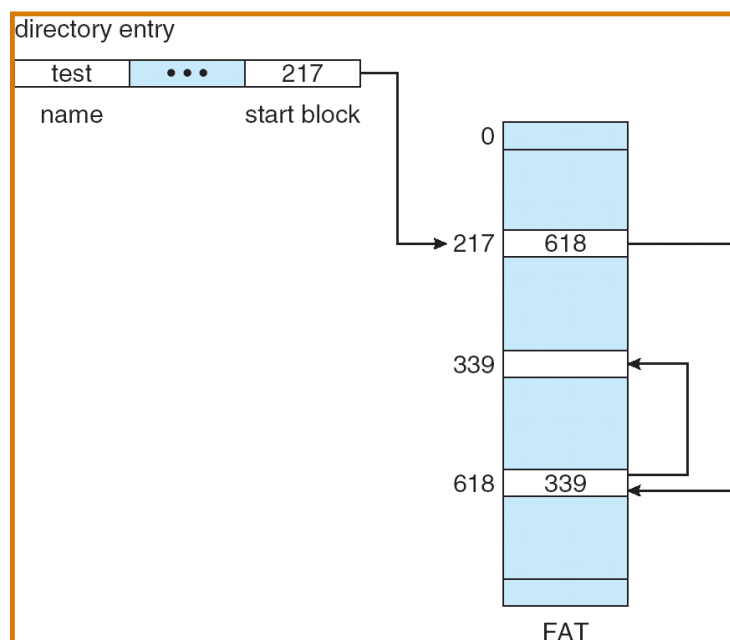


- + Simple - need only starting address
- + Free-space management system - no waste of space
- + Defragmentation not necessary
- No random access
- Extra space required for pointers
- Reliability: what if a pointer gets corrupted?

## Linked Allocation



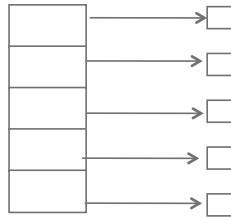
## File-Allocation Table





## Indexed Allocation

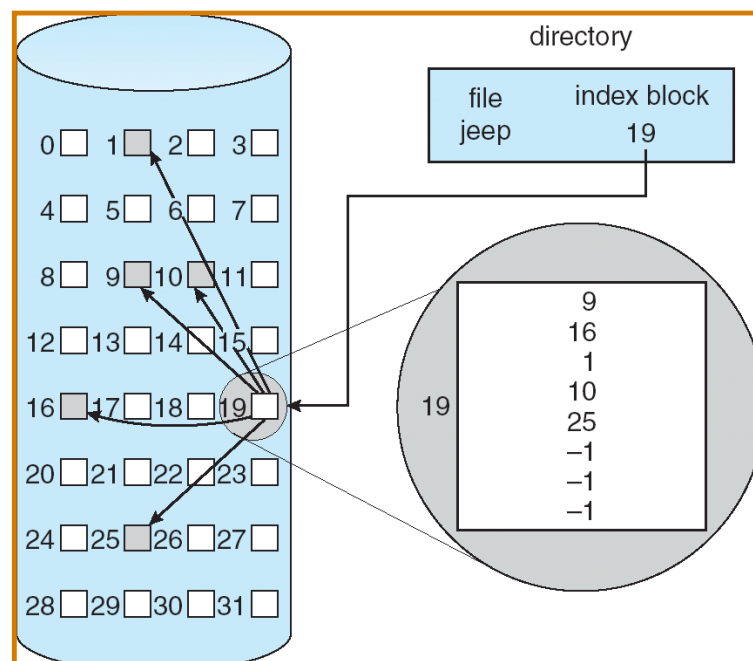
- Brings all pointers together into the *index block*, to allow random access to file blocks.
- Logical view.



index table

- + Supports direct access
- + Prevents external fragmentation
- High pointer overhead --> wasted space

## Example of Indexed Allocation



## Free Space Management

- Disk space limited
- Need to re-use the space from deleted files
- To keep track of free disk space, the system maintains a **free-space list**
  - Records all free disk blocks
- Implemented using
  - Bit vectors
  - Linked lists

## Free-Space Management (Cont.)

- Bit vector ( $n$  blocks)

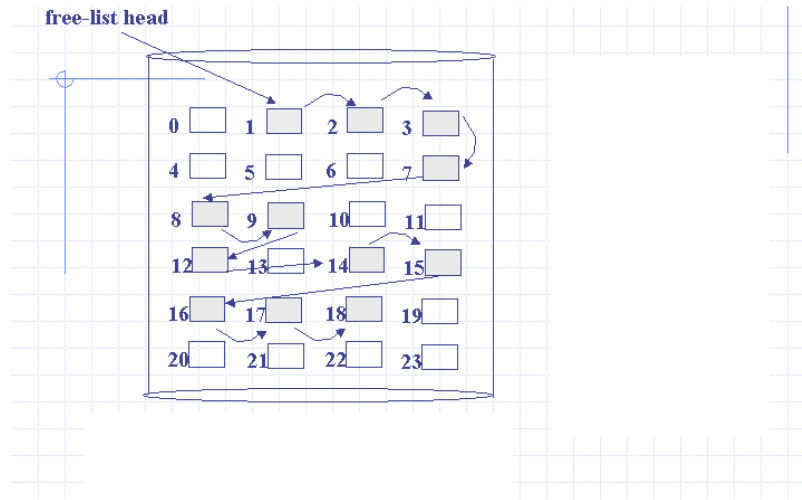


$$\text{bit}[i] = \begin{cases} 0 \Rightarrow \text{block}[i] \text{ free} \\ 1 \Rightarrow \text{block}[i] \text{ occupied} \end{cases}$$

■ e.g. 0000111110001000100010000

## Free-Space Management (Cont.)

- Linked List

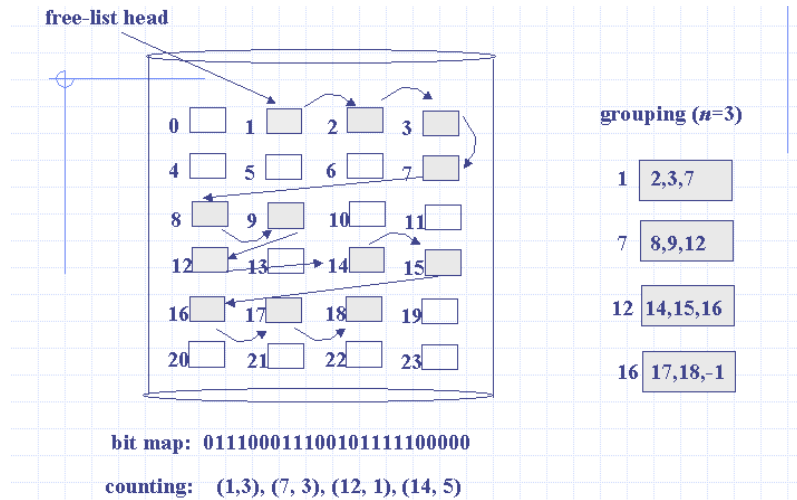


## Free-Space Management (Cont.)

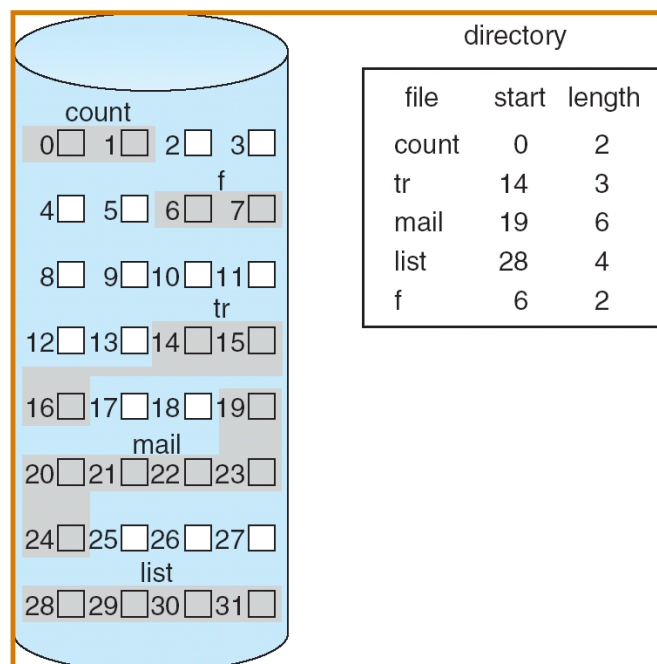
- Bit map requires extra space
  - Example:  
block size =  $2^{12}$  bytes  
disk size =  $2^{30}$  bytes (1 gigabyte)  
 $n = 2^{30}/2^{12} = 2^{18}$  bits (or 32K bytes)
- Easy to get contiguous files
- **Linked list (free list)**
  - Cannot get contiguous space easily
  - requires substantial I/O
- **Grouping**
  - Modification of free-list
  - Store addresses of n free blocks in the first free block
- **Counting**
  - Rather than keeping list of n free addresses:
    - Keep the address of the first free block
    - And the number n of free contiguous blocks that follow it

## Free-Space Management (Cont.)

- Linked List



## Exercise



## Any Questions?



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## Acknowledgements

- “Operating Systems Concepts” book and supplementary material by A. Silberschatz, P. Galvin and G. Gagne
- “Operating Systems: Internals and Design Principles” book and supplementary material by W. Stallings
- “Modern Operating Systems” book and supplementary material by A. Tanenbaum
- R. Doursat and M. Yuksel from UNR

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