CSC 4103 - Operating Systems Fall 2009

LECTURE - XX
FILE SYSTEMS

Tevfik Koşar

Louisiana State University November 5th, 2009

File Systems

- Provides organized and efficient access to data on secondary storage:
 - 1. Organizing data into files and directories and supporting primitives to manipulate them (create, delete, read, write etc)
 - 2. Improve I/O efficiency between disk and memory (perform I/O in units of blocks rather than bytes)
 - 3. Ensure confidentiality and integrity of data
 - Contains file structure via a File Control Block (FCB)
 Ownership, permissions, location..

A Typical File Control Block

file permissions

file dates (create, access, write)

file owner, group, ACL

file size

file data blocks or pointers to file data blocks

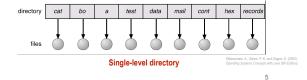
Directories

- > Directories are special files that keep track of other files
 - ✓ the collection of files is systematically organized
 - first, disks are split into partitions that create logical volumes (can be thought of as "virtual disks")
 - \checkmark second, each partition contains information about the files within
 - this information is kept in entries in a device directory (or volume table of contents)
 - ✓ the directory is a symbol table that translates file names into their entries in the directory
 - it has a logical structure
 - it has an implementation structure (linked list, table, etc.)

4

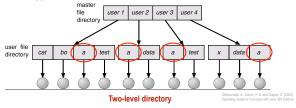
Directories

- Single-level directory structure
 - simplest form of logical organization: one global or root directory containing all the files
 - ✓ problems
 - global namespace: unpractical in multiuser systems
 - no systematic organization, no groups or logical categories of files that belong together



Directories

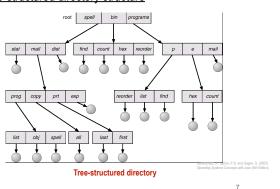
- Two-level directory structure
 - in multiuser systems, the next step is to give each user their own private directory
 - ✓ avoids filename confusion
 - ✓ however, still no grouping: not satisfactory for users with many files



6

Directories

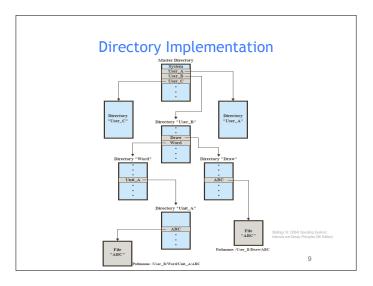
Tree-structured directory structure



Directories

- > Tree-structured directory structure
 - ✓ natural extension of the two-level scheme
 - ✓ provides a general hierarchy, in which files can be grouped in natural ways
 - ✓ good match with human cognitive organization: tendency to categorize objects in embedded sets and subsets
 - ✓ navigation through the tree relies on pathnames
 - absolute pathnames start from the root, example: /jsmith/ academic/teaching/cs446/assignment4/grades
 - relative pathnames start at from a current working directory, example: assignment4/grades
 - the current and parent directory are referred to as . and ..

8



Directory Implementation

- Linear list of file names with pointer to the data blocks.
 - simple to program
 - time-consuming to execute
- Hash Table linear list with hash data structure.
 - decreases directory search time
 - **collisions** situations where two file names hash to the same location
 - fixed size

Allocation Methods

- An allocation method refers to how disk blocks are allocated for files:
- · Contiguous allocation
- Linked allocation
- · Indexed allocation

Contiguous Allocation

- Each file occupies a set of contiguous blocks on the disk
- + Simple only starting location (block #) and length (number of blocks) are required
- Wasteful of space (dynamic storage-allocation problem fragmentation)
- · Files cannot grow

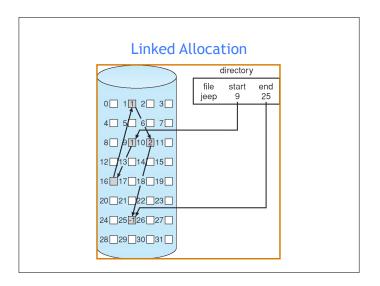
Contiguous Allocation of Disk Space directory file start length count 0 1 2 3 count 0 2 3 tr 14 4 5 6 7 mail 19 6 list 28 4 8 9 10 11 2 12 13 14 15 16 17 18 19 mail 20 21 22 23 24 25 26 27 list 28 29 30 31

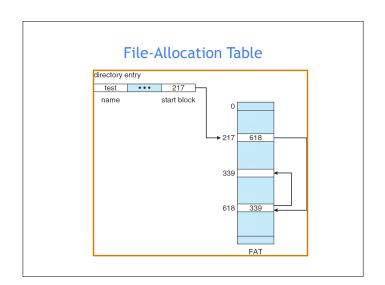
Linked Allocation

• Each file is a linked list of disk blocks: blocks may be scattered anywhere on the disk.



- + Simple need only starting address
- + Free-space management system no waste of space
- + Defragmentation not necessary
- No random access
- Extra space required for pointers
- Reliability: what if a pointer gets corrupted?



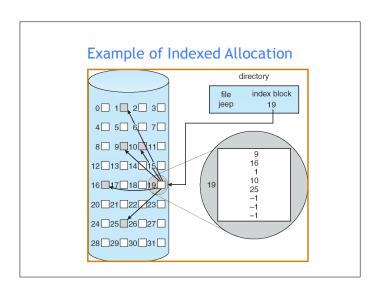


Indexed Allocation

- Brings all pointers together into the *index block*, to allow random access to file blocks.
- · Logical view.



- + Supports direct access
- + Prevents external fragmentation
- High pointer overhead --> wasted space



Free Space Management

- · Disk space limited
- Need to re-use the space from deleted files
- To keep track of free disk space, the system maintains a free-space list
 - Records all free disk blocks
- Implemented using
 - Bit vectors
 - Linked lists

Free-Space Management (Cont.)

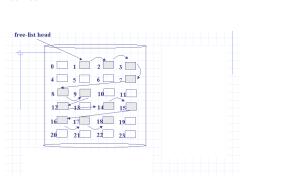
• Bit vector (*n* blocks)



e.g. 0000111110001000100010000

Free-Space Management (Cont.)

Linked List



Free-Space Management (Cont.)

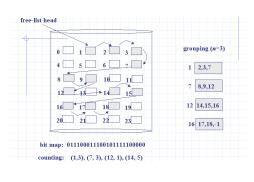
- Bit map requires extra space
 - Example:

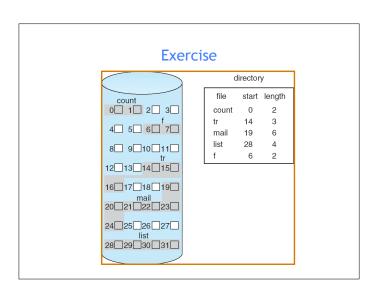
block size = 2^{12} bytes disk size = 2^{30} bytes (1 gigabyte) $n = 2^{30}/2^{12} = 2^{18}$ bits (or 32K bytes)

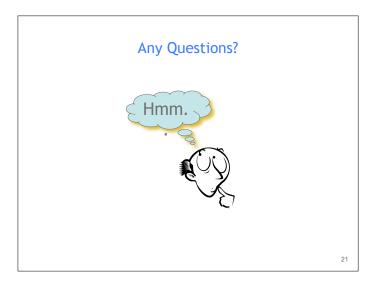
- Easy to get contiguous files
- Linked list (free list)
 - Cannot get contiguous space easily
 - requires substantial I/O
- Grouping
 - Modification of free-list
 - Store addresses of n free blocks in the first free block
- Counting
 - Rather than keeping list of n free addresses:
 - Keep the address of the first free block
 - And the number n of free contiguous blocks that follow it

Free-Space Management (Cont.)

Linked List







Acknowledgements

- "Operating Systems Concepts" book and supplementary material by A. Silberschatz, P. Galvin and G. Gagne
- "Operating Systems: Internals and Design Principles" book and supplementary material by W. Stallings
- "Modern Operating Systems" book and supplementary material by A. Tanenbaum
- R. Doursat and M. Yuksel from UNR

26