

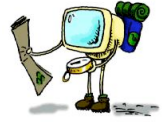
LECTURE - XVI  
MAIN MEMORY - I

Tevfik Koşar

Louisiana State University  
October 22<sup>nd</sup>, 2009

## Roadmap

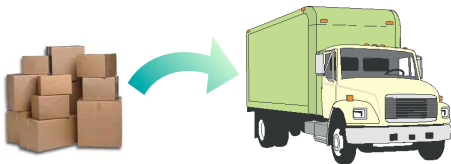
- Midterm Discussion
- Project 1 Discussion
- Main Memory
  - Introduction
  - Contiguous Allocation
  - Dynamic Allocation Algorithms
  - Fragmentation
  - Address Binding
  - Address Protection



2

## Memory Management Requirements

- The O/S must fit multiple processes in memory
  - ✓ memory needs to be subdivided to accommodate multiple processes
  - ✓ memory needs to be allocated to ensure a reasonable supply of ready processes so that the CPU is never idle
  - ✓ memory management is an **optimization** task under **constraints**



Fitting processes into memory is like fitting boxes into a fixed amount of space

3

## Memory Allocation - contiguous

- Fixed-partition allocation
  - Divide memory into fixed-size partitions
  - Each partition contains exactly one process
  - **The degree of multi programming is bound by the number of partitions**
  - When a process terminates, the partition becomes available for other processes

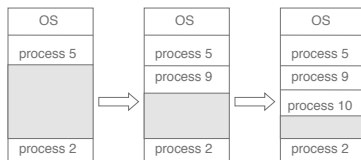
OS
process 5
process 9
process 10
process 2

→ no longer in use

12

## Memory Allocation (Cont.)

- Variable-partition Scheme (Dynamic)
  - When a process arrives, search for a hole large enough for this process
  - Hole - block of available memory; holes of various size are scattered throughout memory
  - Allocate only as much memory as needed
  - Operating system maintains information about:
    - a) allocated partitions    b) free partitions (hole)



13

## Dynamic Storage-Allocation Problem

How to satisfy a request of size  $n$  from a list of free holes

- **First-fit:** Allocate the *first* hole that is big enough
- **Best-fit:** Allocate the *smallest* hole that is big enough; must search entire list, unless ordered by size. Produces the smallest leftover hole.
- **Worst-fit:** Allocate the *largest* hole; must also search entire list. Produces the largest leftover hole.

First-fit and best-fit better than worst-fit in terms of speed and storage utilization

14

## Example

Given five memory partitions of 100 KB, 500 KB, 200 KB, 300 KB, and 600 KB (in order), how would each of the first-fit, best-fit, and worst-fit algorithms place processes of 212 KB, 417 KB, 112 KB, and 426 KB (in order)? Which algorithm makes the most efficient use of memory?

7

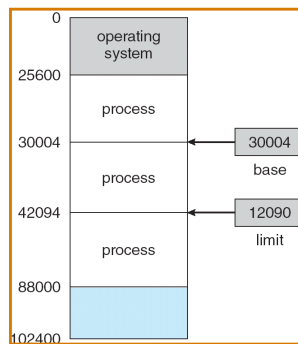
## Fragmentation

- **External Fragmentation** - total memory space exists to satisfy a request, but it is not contiguous (in average ~50% lost)
- **Internal Fragmentation** - allocated memory may be slightly larger than requested memory; this size difference is memory internal to a partition, but not being used
- Reduce external fragmentation by **compaction**
  - Shuffle memory contents to place all free memory together in one large block
  - Compaction is possible *only* if relocation is dynamic, and is done at execution time

15

## Logical Address Space

- Each process has a separate memory space
- Two registers provide address protection between processes:
  - **Base register**: smallest legal address space
  - **Limit register**: size of the legal range



9

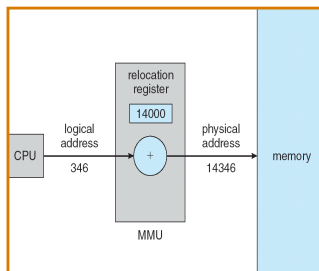
## Address Binding

- Addresses in a source program are generally **symbolic**
  - eg. int count;
- A compiler **binds** these symbolic addresses to **relocatable** addresses
  - eg. 100 bytes from the beginning of this module
- The linkage editor or loader will in turn bind the relocatable addresses to **absolute** addresses
  - eg. 74014
- Each binding is **mapping** from one address space to another

10

## Memory-Management Unit (MMU)

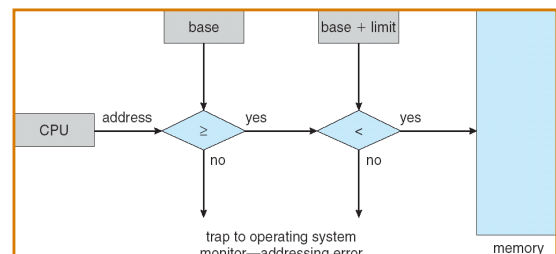
- Hardware device that maps virtual to physical address
- In MMU scheme, the value in the **relocation register** (base register) is added to every address generated by a user process at the time it is sent to memory
- The **user program** deals with **logical** addresses; it **never sees** the **real** physical addresses



11

## HW Address Protection

- CPU hardware compares every address generated in user mode with the registers
- Any attempt to access other processes' memory will be trapped and cause a **fatal error**



12

## Summary

- Main Memory
  - Contiguous Allocation
  - Dynamic Allocation Algorithms
  - Fragmentation
  - Address Binding
  - Address Protection



- Next Lecture: Continue with Main Memory
- Reading Assignment: Chapter 8 from Silberschatz.

13

## Acknowledgements

- “Operating Systems Concepts” book and supplementary material by A. Silberschatz, P. Galvin and G. Gagne
- “Operating Systems: Internals and Design Principles” book and supplementary material by W. Stallings
- “Modern Operating Systems” book and supplementary material by A. Tanenbaum
- R. Doursat and M. Yuksel from UNR

14