

LECTURE - XXI  
PROTECTION AND SECURITY - II

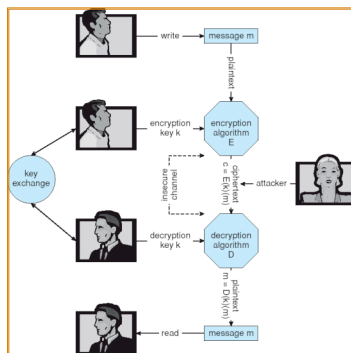
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## Symmetric Encryption

- Same key used to encrypt and decrypt
  - $E(k)$  can be derived from  $D(k)$ , and vice versa
- DES is most commonly used symmetric block-encryption algorithm (created by US Govt)
  - Encrypts a block of data at a time (64 bit messages, with 56 bit key)
- Triple-DES considered more secure (repeat DES three times with three different keys)
- Advanced Encryption Standard (AES) replaces DES
  - Key length upto 256 bits, working on 128 bit blocks
- Twofish, RC4, RC5 .. other symmetric algorithms
- RC4 is most common symmetric stream cipher (works on bits, not blocks), but known to have vulnerabilities
  - Encrypts/decrypts a stream of bytes (i.e wireless transmission, web browsers)
  - Key is a input to pseudo-random-bit generator
    - Generates an infinite **keystream**

## Symmetric Encryption



## Asymmetric Encryption

- Encryption and decryption keys are different
- Public-key encryption based on each user having two keys:
  - public key - published key used to encrypt data
  - private key - key known only to individual user used to decrypt data
- Must be an encryption scheme that can be made public without making it easy to figure out the decryption scheme
  - Most common is RSA (Rivest, Shamir, Adleman) block cipher

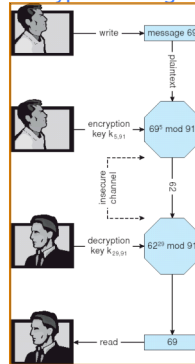
## Asymmetric Encryption (Cont.)

- Formally, it is computationally infeasible to derive  $D(k_d, N)$  from  $E(k_e, N)$ , and so  $E(k_e, N)$  need not be kept secret and can be widely disseminated
  - $E(k_e, N)$  (or just  $k_e$ ) is the **public key**
  - $D(k_d, N)$  (or just  $k_d$ ) is the **private key**
  - $N$  is the product of two large, randomly chosen prime numbers  $p$  and  $q$  (for example,  $p$  and  $q$  are 512 bits each)
  - Select  $k_e$  and  $k_d$ , where  $k_e$  satisfies  $k_e k_d \text{ mod } (p-1)(q-1) = 1$
  - Encryption algorithm is  $E(k_e, N)(m) = m^{k_e} \text{ mod } N$
  - Decryption algorithm is then  $D(k_d, N)(c) = c^{k_d} \text{ mod } N$

## Asymmetric Encryption Example

- For example, choose  $p = 7$  and  $q = 13$
- We then calculate  $N = 7 \cdot 13 = 91$  and  $(p-1)(q-1) = 72$
- We next select  $k_e$  relatively prime to 72 and  $< 72$ , yielding 5
- Finally, we calculate  $k_d$  such that  $k_e k_d \text{ mod } 72 = 1$ , yielding 29
- We now have our keys
  - Public key,  $k_e, N = 5, 91$
  - Private key,  $k_d, N = 29, 91$
- Encrypting the message 69 with the public key results in the ciphertext 62 ( $E = 69^5 \text{ mod } 91$ )
- Ciphertext can be decoded with the private key
  - Public key can be distributed in plaintext to anyone who wants to communicate with holder of public key

## Encryption and Decryption using RSA Asymmetric



## Cryptography (Cont.)

- Note symmetric cryptography based on transformations, asymmetric based on mathematical functions
  - Asymmetric much more compute intensive
  - Typically not used for bulk data encryption
  - Used for authentication, confidentiality, key distribution

## Authentication

- Constraining set of potential senders of a message
  - Complementary and sometimes redundant to encryption
  - Also can prove message unmodified
- Algorithm components
  - A set  $K$  of keys
  - A set  $M$  of messages
  - A set  $A$  of authenticators
  - A function  $S : K \rightarrow (M \rightarrow A)$ 
    - That is, for each  $k \in K$ ,  $S(k)$  is a function for generating authenticators from messages
    - Both  $S$  and  $S(k)$  for any  $k$  should be efficiently computable functions
  - A function  $V : K \rightarrow (M \times A \rightarrow \{\text{true}, \text{false}\})$ . That is, for each  $k \in K$ ,  $V(k)$  is a function for verifying authenticators on messages
    - Both  $V$  and  $V(k)$  for any  $k$  should be efficiently computable functions

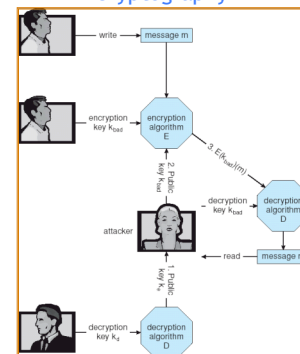
## Authentication (Cont.)

- For a message  $m$ , a computer can generate an authenticator  $a \in A$  such that  $V(k)(m, a) = \text{true}$  only if it possesses  $S(k)$
- Thus, computer holding  $S(k)$  can generate authenticators on messages so that any other computer possessing  $V(k)$  can verify them
- Computer not holding  $S(k)$  cannot generate authenticators on messages that can be verified using  $V(k)$
- Since authenticators are generally exposed (for example, they are sent on the network with the messages themselves), it must not be feasible to derive  $S(k)$  from the authenticators

## Key Distribution

- Delivery of symmetric key is huge challenge
  - Sometimes done **out-of-band**, via paper documents or conversation
- Asymmetric keys can proliferate - stored on **key ring**
  - Even asymmetric key distribution needs care - man-in-the-middle attack

## Man-in-the-middle Attack on Asymmetric Cryptography



## Digital Certificates

- Proof of who or what owns a public key
- Public key digitally signed a trusted party
- Trusted party receives proof of identification from entity and certifies that public key belongs to entity
- Certificate authority are trusted party - their public keys included with web browser distributions
  - They vouch for other authorities via digitally signing their keys, and so on

## Encryption Example - SSL

- Insertion of cryptography at one layer of the ISO network model (the transport layer)
- SSL - Secure Socket Layer (also called TLS)
- Cryptographic protocol that limits two computers to only exchange messages with each other
  - Very complicated, with many variations
- Used between web servers and browsers for secure communication (credit card numbers)
- The server is verified with a **certificate** assuring client is talking to correct server
- Asymmetric cryptography used to establish a secure **session key** (symmetric encryption) for bulk of communication during session
- Communication between each computer then uses symmetric key cryptography

## User Authentication

- Crucial to identify user correctly, as protection systems depend on user ID
- User identity most often established through *passwords*, can be considered a special case of either keys or capabilities
  - Also can include something user has and /or a user attribute
- A password can be associated with each resource (eg. File)
- Different passwords may be associated with different access rights
  - Eg. Reading, updating, and deleting files
- Passwords must be kept secret
  - Frequent change of passwords
  - Use of "non-guessable" passwords
  - Log all invalid access attempts
- Passwords may also either be encrypted or allowed to be used only once

## Password Vulnerabilities

- Password length
  - A four digit password would take less than 5 seconds to crack
- Password combination
  - Should use combination of digits, upper and lower case letters, and other characters
- Never write your password somewhere, memorize it
- Periodically change your password
- Do not use the following in your password:
  - Name, lastname
  - Username
  - Date of birth, zipcode, other personal info
- Do not share your accounts with others

## Encrypted Passwords

- How to keep a password secure within the computer?
- UNIX-type systems keep the password lists encrypted
  - Impossible to invert
  - Simple to compute

=> one-way encryption
- Comparison is performed between encoded passwords
- Another level of protection:
  - Encrypted password file is only readable to root

## Biometrics

- Instead of passwords, use biometric measures
  - Palm-readers
  - Finger-print-readers
  - Iris scanners
  - Voice recognition
- Multi-factor authentication
  - Use a combination of different authentication mechanisms

## Implementing Security Defenses

- **Defense in depth** is most common security theory: using multiple layers of security
- Security policies
  - Eg. Policies on user passwords and accounts
- Vulnerability assessment compares real state of system / network compared to security policy
  - Eg. Assessment to passwords, network ports
- Intrusion detection endeavors to detect attempted or successful intrusions
  - **Signature-based detection**
    - Examine system input or network traffic for specific behavior patterns
  - **Anomaly detection**
    - Detect differences from normal behavior
    - Can also detect previously unknown methods of intrusion: **zero-day attacks**
  - **False-positives** (false alarms) and **false-negatives** (mussed intrusions) are problem
- Auditing, accounting, and logging of all or specific system or network activities

## Any Questions?



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## Reading Assignment

- Read chapter 14 and 15 from Silberschatz.

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## Acknowledgements

- “Operating Systems Concepts” book and supplementary material by Silberschatz, Galvin and Gagne.

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