CSC 4103 - Operating Systems
Spring 2007

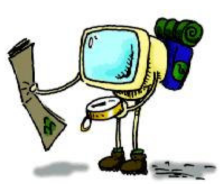
Lecture - IV
- Processes & Threads

Tevfik Koşar

Louisiana State University
January 25th, 2007

Roadmap

- Processes
  - Process Termination
  - Producer-Consumer Problem
  - Inter-process Communication
- Threads
  - Threads vs Processes
  - Multi-threading Models
  - Threading Issues



Process Termination

- Process executes last statement and asks the operating system to delete it (**exit**)
  - Output data from child to parent (via **wait**)
  - Process' resources are deallocated by operating system
- Parent may terminate execution of children processes (**abort**)
  - Child has exceeded allocated resources
  - Task assigned to child is no longer required
  - If parent is exiting
    - Some operating system do not allow child to continue if its parent terminates
      - All children terminated - *cascading termination*

3

Cooperating Processes

- **Independent** process cannot affect or be affected by the execution of another process
- **Cooperating** process can affect or be affected by the execution of another process
- Advantages of process cooperation
  - Information sharing
  - Computation speed-up
  - Modularity
  - Convenience

4

Producer-Consumer Problem

- Paradigm for cooperating processes, *producer* process produces information that is consumed by a *consumer* process
  - *unbounded-buffer* places no practical limit on the size of the buffer
  - *bounded-buffer* assumes that there is a fixed buffer size

5

Bounded-Buffer – Shared-Memory Solution

- Shared data

```
#define BUFFER_SIZE 10
Typedef struct {
   . . .
} item;

item buffer[BUFFER_SIZE];
int in = 0;
int out = 0;
```

- Solution is correct, but can only use BUFFER_SIZE-1 elements

6

Bounded-Buffer – Insert() Method

```
while (true) {
  /* Produce an item */
   while (((in = (in + 1) % BUFFER SIZE
count)  == out)
    ;   /* do nothing -- no free buffers */
   buffer[in] = item;
   in = (in + 1) % BUFFER SIZE;
 {
```

7

Bounded Buffer – Remove() Method

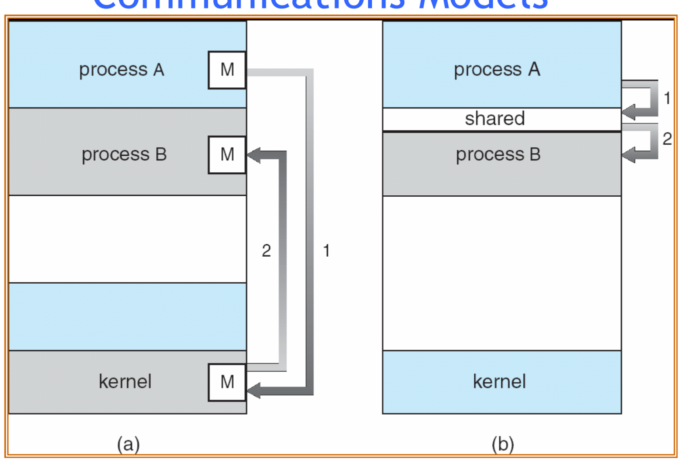
```
while (true) {
    while (in == out)
         ; // do nothing -- nothing to
consume

   // remove an item from the buffer
   item = buffer[out];
   out = (out + 1) % BUFFER SIZE;
return item;
 {
```

8

Interprocess Communication (IPC)

- Mechanism for processes to communicate and to synchronize their actions
- Shared Memory: by using the same address space and shared variables
- Message Passing: processes communicate with each other without resorting to shared variables
- Message Passing facility provides two operations:
  - **send**(*message*) – message size fixed or variable
  - **receive**(*message*)
- If *P* and *Q* wish to communicate, they need to:
  - establish a *communication link* between them
  - exchange messages via send/receive

Communications Models



a) Message Passing b) Shared Memory

Message Passing – direct communication

- Processes must name each other explicitly:
  - **send** (*P*, *message*) – send a message to process P
  - **receive**(*Q*, *message*) – receive a message from process Q
- Properties of communication link
  - Links are established automatically
  - A link is associated with exactly one pair of communicating processes
  - Between each pair there exists exactly one link
  - The link may be unidirectional, but is usually bi-directional
- Symmetrical vs Asymmetrical direct communication
  - **send** (*P*, *message*) – send a message to process P
  - **receive**(id, *message*) – receive a message from any process
- Disadvantage of both: limited modularity, hardcoded

11

Message Passing - indirect communication

- Messages are directed and received from mailboxes (also referred to as ports)
  - Each mailbox has a unique id
  - Processes can communicate only if they share a mailbox
- Primitives are defined as:

  send(*A*, *message*) – send a message to mailbox A

  receive(*A*, *message*) – receive a message from mailbox A

12

Indirect Communication *(cont.)*

- Operations
  - create a new mailbox
  - send and receive messages through mailbox
  - destroy a mailbox
- Properties of communication link
  - Link established only if processes share a common mailbox
  - A link may be associated with many processes
  - Each pair of processes may share several communication links
  - Link may be unidirectional or bi-directional

13

Indirect Communication *(cont.)*

- Mailbox sharing
  - P_1, P_2, and P_3 share mailbox A
  - P_1, sends; P_2 and P_3 receive
  - Who gets the message?
- Solutions
  - Allow a link to be associated with at most two processes
  - Allow only one process at a time to execute a receive operation
  - Allow the system to select arbitrarily the receiver. Sender is notified who the receiver was.

14

Synchronization

- Message passing may be either blocking or non-blocking
- **Blocking** is considered **synchronous**
  - **Blocking send** has the sender block until the message is received
  - **Blocking receive** has the receiver block until a message is available
- **Non-blocking** is considered **asynchronous**
  - **Non-blocking** send has the sender send the message and continue
  - **Non-blocking** receive has the receiver receive a valid message or null

15

Buffering

- Queue of messages attached to the link; implemented in one of three ways
  1. Zero capacity – 0 messages
     Sender must wait for receiver (rendezvous)
  2. Bounded capacity – finite length of *n* messages
     Sender must wait if link full
  3. Unbounded capacity – infinite length
     Sender never waits

16

Threads

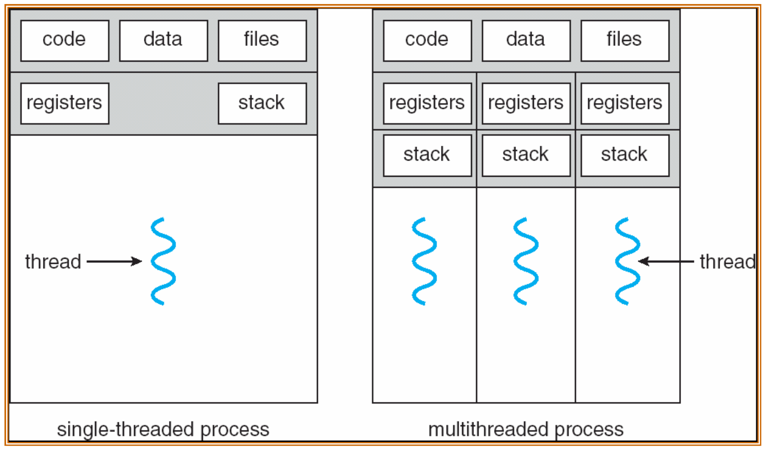
17

Motivation

- In certain cases, a single application may need to run several tasks at the same time
  - Create a new process for each task
    - Time consuming
  - Use a single process with multiple threads

18

Single and Multithreaded Processes

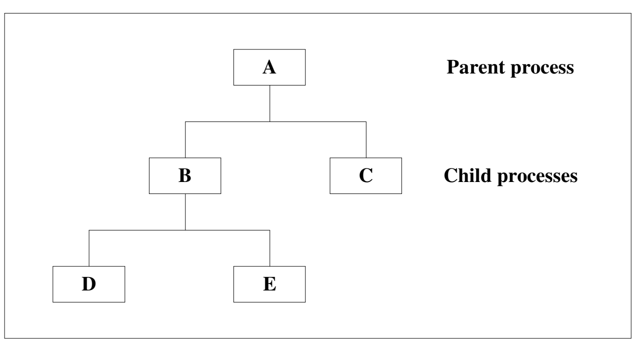


19

Multi-process model

Process Spawning:

Process creation involves the following four main actions:
- **setting up the process control block,**
- **allocation of an address space and**
- **loading the program into the allocated address space and**
- **passing on the process control block to the scheduler**

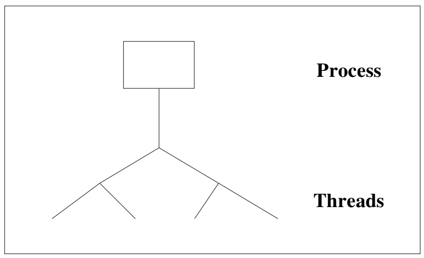


20

Multi-thread model

Thread Spawning:

• **Threads are created *within and belonging to* processes**
• **All the threads created within one process share the resources of the process including the address space**
• **Scheduling is performed on a per-thread basis.**
• **The thread model is a *finer grain scheduling model* than the process model**
• **Threads have a similar *lifecycle* as the processes and will be managed mainly in the same way as processes are**



Threads vs Processes

- Heavyweight Process = Process
- Lightweight Process = Thread

Advantages (Thread vs. Process):
- Much quicker to create a thread than a process
- Much quicker to switch between threads than to switch between processes
- Threads share data easily

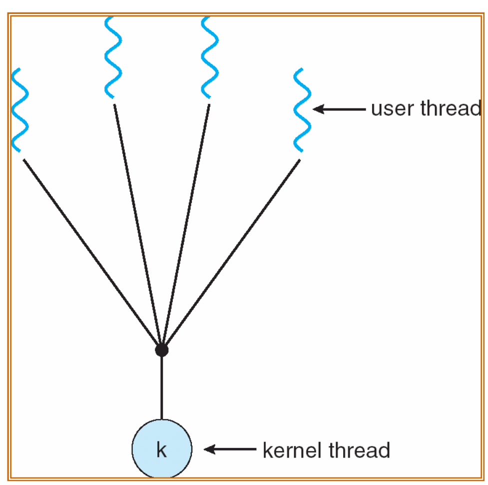
Disadvantages (Thread vs. Process):
- Processes are more flexible
  - They don’t have to run on the same processor
- No security between threads: One thread can stomp on another thread's data
- For threads which are supported by user thread package instead of the kernel:
  - If one thread blocks, all threads in task block.

Different Multi-threading Models

- Many-to-One
- One-to-One
- Many-to-Many

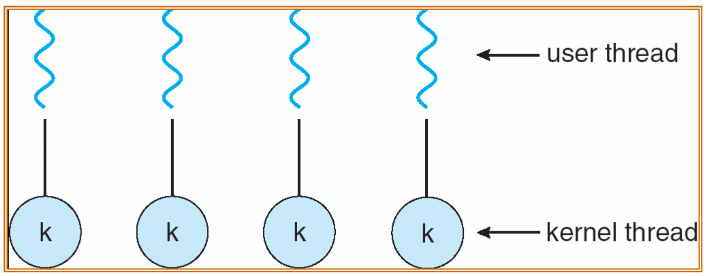
Many-to-One Model

- Several user-level threads mapped to single kernel thread
- Thread management in user space → efficient
- If a thread blocks, entire process blocks
- One thread can access the kernel at a time → limits parallelism
- Examples:
  - Solaris Green Threads
  - GNU Portable Threads



One-to-One Model

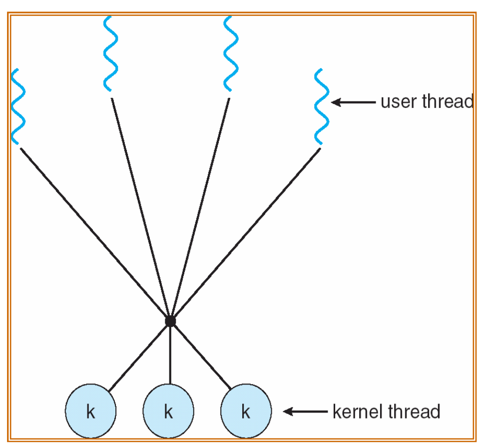
- Each user-level thread maps to a kernel thread
- A blocking thread does not block other threads
- Multiple threads can access kernel concurrently → increased parallelism
- Drawback: Creating a user level thread requires creating a kernel level thread → increased overhead and limited number of threads
- Examples: Windows NT/XP/2000, Linux, Solaris 9 and later



25

Many-to-Many Model

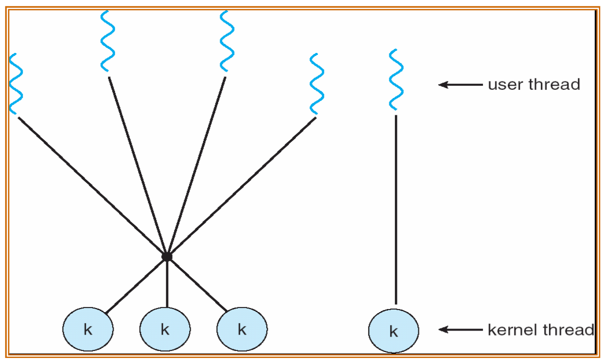
- Allows many user level threads to be mapped to a smaller number of kernel threads
- Allows the operating system to create a sufficient number of kernel threads
- Increased parallelism as well as efficiency
- Solaris prior to version 9
- Windows NT/2000 with the *ThreadFiber* package



26

Two-level Model

- Similar to M:M, except that it allows a user thread to be **bound** to kernel thread
- Examples: IRIX, HP-UX, Tru64 UNIX, Solaris 8 and earlier



27

Threading Issues

- Semantics of **fork()** and **exec()** system calls
- Thread cancellation
- Signal handling
- Thread pools
- Thread specific data

28

Semantics of fork() and exec()

- Semantics of **fork()** and **exec()** system calls change in a multithreaded program
  - Eg. if one thread in a multithreaded program calls fork()
    - Should the new process duplicate all threads?
    - Or should it be single-threaded?
  - Some UNIX systems implement two versions of fork()

  - If a thread executes exec() system call
    - Entire process will be replaced, including all threads

29

Thread Cancellation

- Terminating a thread before it has finished
  - If one thread finishes the searching a database, others may be terminated
  - If user presses a button on a web browser, web page can be stopped from loading further
- Two approaches to cancel the target thread
  - **Asynchronous cancellation** terminates the target thread immediately
  - **Deferred cancellation** allows the target thread to periodically check if it should be cancelled
    - More controlled and safe

30

Signal Handling

- Signals are used in UNIX systems to notify a process that a particular event has occurred
- All signals follow this pattern:
  1. Signal is generated by particular event
  2. Signal is delivered to a process
  3. Once delivered, a signal must be handled
- In multithreaded systems, there are 4 options:
  - Deliver the signal to the thread to which the signal applies
  - Deliver the signal to every thread in the process
  - Deliver the signal to certain threads in the process
  - Assign a specific thread to receive all signals for the process

31

Thread Pools

- Threads come with some overhead as well
- Unlimited threads can exhaust system resources, such as CPU or memory
- Create a number of threads at process startup) and put them in a pool, where they await work
- When a server receives a request, it awakens a thread from this pool
- Advantages:
  - Usually faster to service a request with an existing thread than create a new thread
  - Allows the number of threads in the application(s) to be bound to the size of the pool
- Number of threads in the pool can be setup according to:
  - Number of CPUs, memory, expected number of concurrent requests

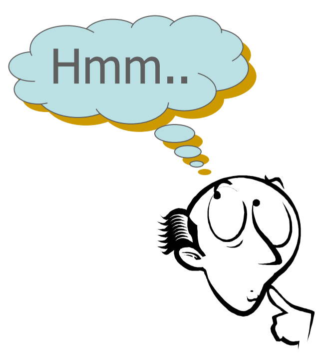
32

Thread Specific Data

- Threads belonging to the same process share the data of the process
- In some cases, each thread needs to have its own copy of data → thread specific
- Useful when you do not have control over the thread creation process (i.e., when using a thread pool)

33

Any Questions?



34

Reading Assignment

- Read chapter 4 from Silberschatz.

35

Acknowledgements

- "Operating Systems Concepts" book and supplementary material by Silberschatz, Galvin and Gagne.

36