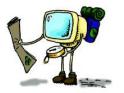
# **Programming Languages**

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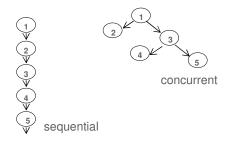
# Roadmap

- Concurrent Programming
  - Shared Memory vs Message Passing
  - Divide and Compute
  - Threads vs Processes
  - Synchronization



## **Concurrent Programming**

- So far, we have focused on sequential programming: all computational tasks are executed in sequence, one after the other.
- Next three lectures, we will focus on concurrent programming: multiple computational tasks are executed simultaneously, at the same time.



## **Concurrent Programming**

- Implementation of concurrent tasks:
  - as separate programs
  - as a set of processes or threads created by a single program
- Execution of concurrent tasks:
  - on a single processor
  - → Multithreaded programming
  - on several processors in close proximity
  - → Parallel computing
  - on several processors distributed across a network
  - → Distributed computing

## **Communication Between Tasks**

Interaction or communication between concurrent tasks can done via:

- Shared memory:
  - all tasks has access to the same physical memory
  - they can communicate by altering the contents of shared memory
- Message passing:
  - no common/shared physical memory
  - tasks communicate by exchanging messages

5

## Motivation

- Increase the performance by running more than one tasks at a time.
  - divide the program to n smaller pieces, and run it n times faster using n processors
- To cope with independent physical devices.
  - do not wait for a blocked device, perform other operations at the background

# Divide and Compute

```
x1 + x2 + x3 + x4 + x5 + x6 + x7 + x8

How many operations with sequential programming?

7

Step 1: x1 + x2
Step 2: x1 + x2 + x3
Step 3: x1 + x2 + x3 + x4
Step 4: x1 + x2 + x3 + x4 + x5
Step 5: x1 + x2 + x3 + x4 + x5 + x6
Step 6: x1 + x2 + x3 + x4 + x5 + x6
Step 7: x1 + x2 + x3 + x4 + x5 + x6 + x7
Step 7: x1 + x2 + x3 + x4 + x5 + x6 + x7 + x8
```

7

## Divide and Compute

```
Step 1: 1 + 2 = 3

Step 2: 3 + 3 = 6

Step 3: 6 + 4 = 10

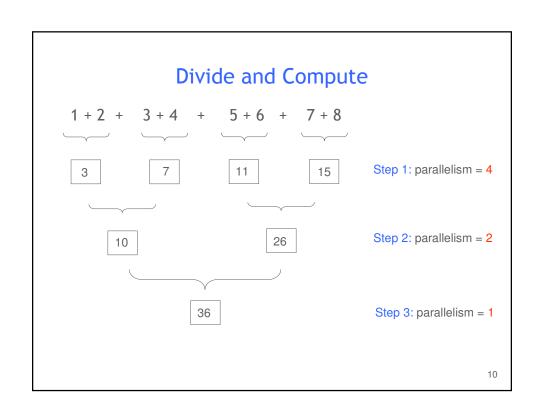
Step 4: 10 + 5 = 15

Step 5: 15 + 6 = 21

Step 6: 21 + 7 = 28

Step 7: 28 + 8 = 36
```

1 + 2 + 3 + 4 + 5 + 6 + 7 + 8



# Gain from parallelism

#### In theory:

 dividing a program into n smaller parts and running on n processors results in n time speedup

### In practice:

- This is not true, due to
  - Communication costs
  - Dependencies between different program parts
    - Eg. the addition example can run only in log(n) time not 1/n

1

## **Prevent Blocking**

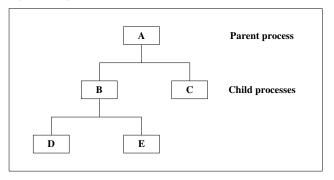
- Do not wait for a blocked device, perform other operations at the background
  - During I/O perform computation
  - During continuous visualization, handle key strokes and I/O
    - Eg. video games
  - While listening to network, perform other operations
    - Listening to multiple sockets at the same time
  - Concurrent I/O, concurrent transfers
    - Eg. Web browsers

### Threads vs Processes

#### **Process Spawning:**

Process creation involves the following four main actions:

- · setting up the process control block,
- allocation of an address space and
- · loading the program into the allocated address space and
- · passing on the process control block to the scheduler

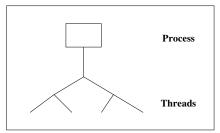


13

### Threads vs Processes

#### Thread Spawning:

- Threads are created within and belonging to processes
- All the threads created within one process share the resources of the process including the address space
- · Scheduling is performed on a per-thread basis.
- The thread model is a *finer grain scheduling model* than the process model
- Threads have a similar *lifecycle* as the processes and will be managed mainly in the same way as processes are



### Threads vs Processes

- Heavyweight Process = Process
- Lightweight Process = Thread

#### Advantages (Thread vs. Process):

- Much quicker to create a thread than a process
- Much quicker to switch between threads than to switch between processes
- Threads share data easily

#### Disadvantages (Thread vs. Process):

- Processes are more flexible
  - They don't have to run on the same processor
- No security between threads: One thread can stomp on another thread's data
- For threads which are supported by user thread package instead of the kernel:
  - If one thread blocks, all threads in task block.

15

# **Synchronization**

 Mechanism that allows the programmer to control the relative order in which operations occur in different threads or processes.

# Synchronization - Threads

Int sum = 0;

Thread 1:
int t;
lock(sum);
sum = sum + x;
t = sum;
....
unlock(sum);

Thread 2:
int t;
lock(sum);
sum = sum + y;
t = sum;
...
unlock(sum);

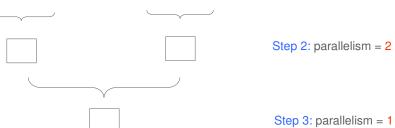
Use of semaphores for thread synchronization!

1

# **Synchronization - Processes**

$$x1 + x2 + x3 + x4 + x5 + x6 + x7 + x8$$

Step 1: parallelism = 4



Wait for a message from other processes before continuing processing!

# On a single processor machine

- You can have multiple threads
- You can also have multiple processes and have the effect of concurrency
  - timesharing

19

#### Data Parallelism TASK 1 a b c for ( i = 1; i<= 20;i++) for (i = 1; i<=60; i++) a(i) = b (i) \* c (i) a(i) = b(i) \* c(i);20 TASK 2 21 } **for (i** = 21; i<= 40; i++) a(i) = b(i) \* c(i);40 41 Decompose DATA into pieces. All the tasks perform the TASK 3 same computation operating with their own piece of data. for (i = 41; i<= 60; i++) a(i) = b(i) \* c(i);20

