

Programming Languages

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Lecture - XIV
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Quiz - 2

1) Consider the following Scheme definition:

```
>> (define (test x y) (if (= x 3) 1 y))
```

Write the output of the following call to this function

```
>> (test 3 (/ 5 0))
```

a) assuming normal-order evaluation is used

b) assuming applicative-order evaluation is used

Quiz - 2

2) Write the output of the following programs:

a) `>> (let ((a 3) (b 4))
 (let ((a 5) (b a))
 (let ((a b) (c a))
 (+ a b c))))`

b) `>> (let ((a 3) (b 4))
 (let* ((a 5) (b a))
 (let* ((a b) (c a))
 (+ a b c))))`

Roadmap

- Scheme
 - List Operations
 - Assignment
 - Loops
 - Evaluation Order in Scheme
- Lambda Calculus
- Comments on Midterm Exam



List Operations

Scheme provides a variety of procedures for operating on lists:

- **length** takes one argument (a list) and returns an integer giving the length of the list.

```
>> (length '(0 #t #f))  
3
```

- **list** takes one or more arguments and constructs a list of those items.

```
>> (list 1)                ; == (cons 1 `())  
(1)  
>> (list 1 2 3)  
(1 2 3)
```

List Operations

- `append` takes two or more lists and constructs a new list with all of their elements.

```
>> (append '(1 2) '(3 4))  
(1 2 3 4)
```

➔ Notice that this is different from what `list` does:

```
>> (list '(1 2) '(3 4))  
((1 2) (3 4))
```

```
>> (append '((1 2) (3 4)) '((5 6) (7 8)))  
((1 2) (3 4) (5 6) (7 8))
```

List Operations

- `reverse` takes one list, and returns a new list with the same elements in the opposite order.

```
>> reverse '(1 2 3 4))  
(4 3 2 1)
```

Assignment

- set!
- set-car!
- set-cdr!

```
>> (let ((x 2)
        (l `(a b)))
      (set x! 3)
      (set-car! l `(c d))
      (set-cdr! l `(e)))
```

```
>> x
```

```
3
```

```
>> l
```

```
((c d) e)
```


Loops

```
>> (do ((i 0 (+ i 1))  
      (a 0 b)  
      (b 1 (+ a b)))  
      ((= i n)) b)  
      (display b)  
      (display " "))))
```

-> initialize & increment

-> “ “

-> “ “

-> termination test

-> body of the loop

-> “ “

Loops

```
>> (for-each (lambda (a b) (display (* a b)) (newline))  
      `(2 4 6)  
      `(3 5 7))
```

6

20

42

()

Evaluation Order in Scheme

- Scheme uses **applicative-order** evaluation
 - Evaluate arguments before passing them to subroutine

```
>> (define double (lambda (x) (+ x x)))
```

Applicative order evaluation (as in Scheme):

```
>> (double (* 3 4))
```

```
→ (double 12)
```

```
→ (+ 12 12)
```

```
→ 24
```

Evaluation Order in Scheme

```
>> (define double (lambda (x) (+ x x)))
```

Normal order evaluation:

```
>> (double (* 3 4))
```

```
→ (+ (* 3 4) (* 3 4))
```

```
→ (+ 12 (* 3 4))
```

```
→ (+ 12 12)
```

```
→ 24
```

→ Normal order causes us to evaluate `(* 3 4)` twice

Evaluation Order in Scheme

- In specific cases, the outcome can be different.

Eg. (define switch (lambda (x a b c)
 (cond ((< x 0) a)
 ((= x 0) b)
 ((> x 0) c))))

Using applicative order:

>> (switch -1 (+ 1 2) (+ 2 3) (+ 3 4))

→ (switch -1 3 (+ 2 3) (+ 3 4))

→ (switch -1 3 5 (+ 3 4))

→ (switch -1 3 5 7)

→ (cond ((< x 0) 3)

((= x 0) 5)

((> x 0) 7))

→ (cond (#t 3)

((= x 0) 5)

((> x 0) 7))

Evaluation Order in Scheme

Using normal order:

```
>> (switch -1 (+ 1 2) (+ 2 3) (+ 3 4))
```

```
→ (switch -1 (+ 1 2) (+ 2 3) (+ 3 4))
```

```
→ (cond      ((< x 0) (+ 1 2))  
              ((= x 0) (+ 2 3))  
              ((> x 0) (+ 3 4)))
```

```
→ (cond      (#t (+ 1 2))  
              ((= x 0) (+ 2 3))  
              ((> x 0) (+ 3 4)))
```

```
→ (+ 1 2)
```

```
→ 3
```

→ Normal order avoids evaluating (+ 2 3) and (+ 3 4)

Comments in Scheme

You can and should put comments in your Scheme programs. **Start a comment with a semicolon**. Scheme will ignore any characters after that on a line. (This is like the `//` comments in C++.)

```
>> (define foo 22) ; define foo with an initial value of 22
```

Lambda Calculus

- lambda calculus is a formal system designed to investigate function definition, function application, and recursion.
- was introduced in 1930's
- the smallest universal programming language
- lambda calculus consists of a single transformation rule (variable substitution) and a single function definition scheme
- lambda calculus is universal in the sense that any computable function can be expressed and evaluated using this formalism

Lambda Calculus

- In lambda calculus, every **expression** stands for a function with a single argument
- the **argument** of the function is in turn a function with a single argument
- and the **value** of the function is another function with a single argument

Eg.

$f(x) = x + 2$ would be expressed as $\lambda x. x + 2$

and the number $f(3)$ would be written as $(\lambda x. x + 2) 3$

Lambda Calculus

- A function of **two variables** is expressed in lambda calculus as a function of one argument which returns a function of one argument

Eg. the function $f(x, y) = x - y$ would be written as:

$\lambda x. \lambda y. x - y.$

A common convention is to abbreviate curried functions as, for instance, $\lambda x y. x - y.$

Example:

$$\begin{aligned} & (\lambda x y. x - y) 7 2 \\ &= (\lambda y. 7 - y) 2 \\ &= 7 - 2 \\ &= 5 \end{aligned}$$

Lambda Calculus

- Consider a function which takes another function as argument and applies it to the argument 3: $\lambda f. f\ 3$.
- This latter function could be applied to our earlier "add-two" function as follows: $(\lambda f. f\ 3)\ (\lambda x. x+2)$.

$$\begin{aligned}\rightarrow & (\lambda f. f\ 3)\ (\lambda x. x + 2) \\ &= (\lambda x. x + 2)\ 3 \\ &= 3 + 2 \\ &= 5\end{aligned}$$

Lambda Calculus

Formal definition:

The set of all lambda expressions can then be described by the following context-free grammar in BNF:

expr	→ identifier
expr	→ λ identifier . expr
expr	→ expr expr

Lambda Calculus

Mapping Lambda Calculus to Scheme:

Math: $f(x) = x + 2$

Lambda Calculus: $\lambda x. x + 2$

Scheme: `(define f (lambda (x) (+ x 2)))`

Math: $f(x, y) = x - y$

Lambda Calculus: $\lambda x. \lambda y. x - y$

Scheme: `(define f (lambda (x y) (- x y)))`

Midterm will cover:

- Ch 1.1-1.6
 - Ch 2.1-2.3
 - Ch 3.1-3.3
 - Ch 4.1-4.6
 - Ch 6.1-6.7
 - Ch 10.1-10.5
 - Scheme
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- Check the class notes! Anything not covered in the class will not be asked!
 - Expect similar questions to the homework assignments and quizzes.

Midterm Hints

- Be sure that you know:
 - difference between compilation and interpretation; lexical, syntactic and semantic analysis; linking and binding
 - how to write regular expressions
 - how to generate NFA and DFA from regular expressions
 - how to generate parse trees
 - difference between top-down and bottom-up parsing; LL vs LR grammars
 - difference between stack-based vs heap-based allocation; static vs dynamic scoping
 - how to generate attribute grammars; decorate parse trees; and generate syntax trees
 - expression evaluation orders; applicative vs normal-order evaluation
 - basic scheme functions; lists, searching, and scoping in Scheme